

1. **(T/F)** Let ε denote the machine epsilon for the **Float64** format. If $x \in \mathbf{R}$ is such that $0 < x < \varepsilon$, then x cannot be represented in the **Float64** format.
2. **(T/F)** Machine multiplication is commutative, meaning that $a\hat{*}b = b\hat{*}a$ for any **Float64** point numbers a and b .
3. **(T/F)** If x is a **Float16** and y is a **Float32** number, then the result of $x + y$ is a **Float64** number.
4. **(T/F)** The only polynomial p of degree at most 3 such that $p(-1) = p(0) = p(1) = 1$ is the constant polynomial $p(x) = 1$.
5. **(T/F)** Given $x_0 < x_1 < x_2 < x_3$ and $y_0, y_1, y_2, y_3 \in \mathbf{R}$, the unique polynomial passing through these data points is given by

$$p(x) = \frac{(x-x_1)(x-x_2)(x-x_3)}{(x_0-x_1)(x_0-x_2)(x_0-x_3)}y_0 + \frac{(x-x_0)(x-x_2)(x-x_3)}{(x_1-x_0)(x_1-x_2)(x_1-x_3)}y_1 \\ + \frac{(x-x_0)(x-x_1)(x-x_3)}{(x_2-x_0)(x_2-x_1)(x_2-x_3)}y_2 + \frac{(x-x_0)(x-x_1)(x-x_2)}{(x_3-x_0)(x_3-x_1)(x_3-x_2)}y_3.$$

6. Given $x_0 < \dots < x_n$ and $y_0, \dots, y_n \in \mathbf{R}$, prove that the constant polynomial p that minimizes the expression $\sum_{i=0}^n |y_i - p(x_i)|^2$ is given by

$$p(x) = \frac{1}{n+1} \sum_{i=0}^n y_i.$$

7. **(T/F)** Let $f(x) = \exp(2x)$, and for any $n \in \mathbf{N}$, let $f_n \in \mathcal{P}_n$ denote the polynomial interpolating f at $n+1$ equidistant points $-1 = x_0 < x_1 < \dots < x_n = 1$. Then

$$\lim_{n \rightarrow \infty} \left(\max_{-1 \leq x \leq 1} |f(x) - f_n(x)| \right) = 0.$$

8. **(T/F)** In Julia, if A is a 10 by 10 matrix, then `A[isodd.(1:end), :]` gives the matrix obtained by keeping only the rows with odd indices.
9. **(T/F)** The degree of precision of the Gauss–Legendre quadrature rule with $n+1$ points is equal to $2n$.
10. Describe in words what the following code does.

```
function I_approx(f, a, b, n)
    x = LinRange(a, b, n + 1)
    h = x[2] - x[1]
    return h * sum(f, x[1:n]) .+ h/2)
end
```

In the next questions, we consider the following linear system:

$$\mathbf{Ax} = \mathbf{b}, \quad \mathbf{A} \in \mathbf{R}^{n \times n}, \quad \mathbf{b} \in \mathbf{R}^n. \quad (1)$$

1. Suppose that \mathbf{A} is a 10 by 10 upper triangular matrix, and that \mathbf{b} is a vector of size 10. Complete the following implementation of the backward substitution algorithm:

```
# Suppose that A and b have already been defined  
 $\mathbf{x} = \text{zero}(\mathbf{b})$   
# YOUR CODE BELOW
```

2. How many floating point operations does your implementation require?
3. Suppose that \mathbf{A} is symmetric and positive definite. Describe step by step an efficient direct method for solving (1) in this case.

4. **(T/F)** Suppose that \mathbf{A} is symmetric and positive definite. Then there exists a unique solution \mathbf{x}_* to (1) and, furthermore,

$$\mathbf{x}_* \in \arg \min_{\mathbf{x} \in \mathbf{R}^n} \left(\frac{1}{2} \mathbf{x}^T \mathbf{A} \mathbf{x} - \mathbf{b}^T \mathbf{x} \right).$$

5. **(T/F)** The Gauss-Seidel basic iterative method is convergent if and only if the matrix \mathbf{A} is strictly diagonally dominant.

Until the end of the quiz, we consider the following scalar, nonlinear equation, where the function f is *twice continuously differentiable*:

$$f(x) = 0, \quad f: \mathbf{R} \rightarrow \mathbf{R}, \quad x \in \mathbf{R}. \quad (2)$$

1. **(T/F)** If $f'(x) > 1$ for all $x \in \mathbf{R}$, then there exists a unique solution to (2).
2. **(T/F)** Suppose that $f'(x) > 1$ for all $x \in \mathbf{R}$ and that the following converges to some $x_* \in \mathbf{R}$ when started from $x_0 = 1$. Then $f(x_*) = 0$.

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)}. \quad (3)$$

3. **(T/F)** If $f(0)f(1) > 0$, then there cannot exist a solution $x_* \in (0, 1)$ to (2).
4. **(2 marks)** Suppose that iteration (3) converges to some $x_* \in \mathbf{R}$. Prove that

$$\lim_{k \rightarrow \infty} \left| \frac{x_{k+1} - x_*}{x_k - x_*} \right| = \left| \frac{f''(x_*)}{2f'(x_*)} \right|.$$